

# Edge degree conditions for subpancyclicity in line graphs

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## Abstract

In this paper, two best possible edge degree conditions are given for the line graph  $L(G)$  of a graph  $G$  with girth at least 4 or 5 to be subpancyclic, i.e.,  $L(G)$  contains a cycle of length  $k$ , for each  $k$  between 3 and the circumference of  $L(G)$ . In [5] the following conjecture is made: If  $G$  is a graph such that the degree sum of any pair of adjacent vertices in  $G$  is greater than  $(\sqrt{8n+1}+1)/2$ , then the line graph  $L(G)$  of  $G$  is pancyclic whenever  $L(G)$  is Hamiltonian, unless  $G$  is isomorphic to  $C_4$ ,  $C_5$ , or the Petersen graph. Our results show that the conjecture is true for those graphs of order  $n \geq 72$  with girth at least 4. © 1998 Elsevier Science B.V. All rights reserved

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## 1. Introduction

We use [1] for terminology and notation not defined here and consider finite simple graphs only. Set  $\lambda(G) = \{k \mid G \text{ contains a cycle of length } k\}$ . A graph  $G$  is *pancyclic* if  $\lambda(G) = [3, |V(G)|]$ , i.e.,  $\lambda(G)$  contains all integers in  $[3, |V(G)|]$ . Similarly, a graph  $G$  is called *subpancyclic* if  $\lambda(G) = [3, \text{cr}(G)]$  where  $\text{cr}(G)$  denotes the *circumference* of a graph  $G$ , i.e. the length of a longest cycle of  $G$ . Let  $g(G)$  denote the *girth* of a graph, i.e., the length of a shortest cycle of  $G$ . The degree of an edge  $e = uv \in E(G)$ , denoted by  $d(e)$ , is defined to be  $d(u) + d(v)$ . The *minimum edge degree* of a graph  $G$ , denoted by  $\delta'(G)$ , is defined to be  $\min\{d(u) + d(v) \mid uv \in E(G)\}$ .

A natural question is the following: how large should be the minimum degree  $\delta(G)$  of a graph  $G$  be in order that the Hamiltonian line graph  $L(G)$  is guaranteed to be pancyclic?

From results in [4] it follows that there exists a constant  $A$  such that if  $L(G)$  is Hamiltonian and  $\delta(G) > A \cdot n^{1/3}$ , then  $L(G)$  is pancyclic. Also, there exists a constant  $B$

and an infinite family of graphs  $G$  such that  $\delta(G) > B \cdot n^{1/3}$  and  $L(G)$  is Hamiltonian but not pancyclic.

Here we consider a similar question concerning edge degree.

Our main results are the following.

**Theorem 1.** *Let  $G$  be a graph of order  $n$  ( $n \geq 72$ ), girth at least 5, and with  $\delta'(G) > \sqrt{2n+1}$  (or  $\delta'^2(G) - 1 > 2n$ ). Then  $L(G)$  is subpancyclic and the result is best possible.*

**Theorem 2.** *Let  $G$  be a graph of order  $n$  ( $n \geq 72$ ), girth at least 4, and with  $\delta'^2(G) - \delta'(G) > 2n$ . Then  $L(G)$  is subpancyclic and the result is best possible.*

In general, results involving degree sums are direct derivatives from results involving the minimum degree of the graph. Theorems 1 and 2 show an exception to this rule (compare the results in [4] mentioned above).

The main theorems in the paper deal with graphs with girth at least 5 (Theorem 1) and girth at least 4 (Theorem 2). Since the necessary degree condition for subpancyclicity in Theorem 1 is a little weaker than the degree condition in Theorem 2, an obvious question is if one can expect an even weaker condition for graphs with girth at least  $k$  for some  $k \geq 6$ . But, in fact, the extremal graphs given at the end of the proof of Theorem 1 can be chosen to have arbitrarily large girth. This shows that even for graphs with large girth the degree condition in Theorem 1 is best possible.

## 2. Proofs of the main theorems

Before proving Theorems 1 and 2 we introduce some additional terminology and notation, and state a number of preliminary results.

By a *circuit* of a graph  $G$  we will mean an Eulerian subgraph of  $G$ , i.e., a connected subgraph in which every vertex has even degree. Note that by this definition (the trivial subgraph induced by) a single vertex is also a circuit. If  $\mathcal{C}$  is a circuit of  $G$ , then  $\tilde{E}(\mathcal{C})$  denotes the set of edges of  $G$  incident with at least one vertex of  $\mathcal{C}$ , we write  $\varepsilon(\mathcal{C})$  for  $|E(\mathcal{C})|$  and  $\bar{\varepsilon}(\mathcal{C})$  for  $|\tilde{E}(\mathcal{C})|$ . Let  $\mathcal{C}_m$  denote a cycle of length  $m$ .

Harary and Nash–Williams [3] characterized Hamiltonian line graphs.

**Theorem 3** (Harary and Nash–Williams [3]). *The line graph  $L(G)$  of a graph  $G$  is Hamiltonian if and only if  $G$  contains a circuit  $\mathcal{C}$  such that  $\bar{\varepsilon}(\mathcal{C}) = |E(G)| \geq 3$ .*

From Theorem 3 one easily prove a more general result (see e.g., [2]).

**Theorem 4.** *The line graph  $L(G)$  of a graph  $G$  contains a cycle of length  $k \geq 3$  if and only if  $G$  contains a circuit  $\mathcal{C}$  such that  $\varepsilon(\mathcal{C}) \leq k \leq \bar{\varepsilon}(\mathcal{C})$ .*

A key lemma for our proof of Theorems 1 and 2 is the following.

**Lemma 5.** Let  $G$  be a graph of order  $n$ , girth at least 4, and with minimum edge degree  $\delta' \geq 8$  such that  $L(G)$  contains  $\mathcal{C}_{m+1}$  but not  $\mathcal{C}_m$ . Then

$$m \leq \frac{4n - \delta'}{\delta'}.$$

**Proof.** Let  $G$  satisfy the hypothesis of the lemma. By Theorem 4,  $G$  contains a circuit  $\mathcal{C}$  with  $\varepsilon(\mathcal{C}) \leq m + 1 \leq \bar{\varepsilon}(\mathcal{C})$ . In fact,  $\varepsilon(\mathcal{C}) = m + 1$ , otherwise  $L(G)$  contains  $\mathcal{C}_m$ . By Theorem 4,  $\varepsilon(\mathcal{C}) \geq \Delta(G) + 2 \geq (\delta'(G) + 4)/2 \geq 6$ . Since  $\mathcal{C}$  is a circuit, there exist edge disjoint cycles  $D_1, \dots, D_r$  such that

$$\mathcal{C} = \bigcup_{i=1}^r D_i.$$

We distinguish the following cases.

*Case 1:*  $r = 1$ . Then  $\mathcal{C}$  is a cycle of length  $m + 1$ .

*Case 1.1:*  $\mathcal{C}$  has a chord. Let  $\mathcal{C}'$  be a longest cycle among all cycles that contain exactly one chord of  $\mathcal{C}$  while the remaining edges belong to  $\mathcal{C}$ .

In  $\sum_{e \in E(\mathcal{C}')} d(e)$ , every edge in  $\bar{E}(\mathcal{C}')$  is counted at most four times. Hence,

$$\bar{\varepsilon}(\mathcal{C}') \geq \varepsilon(\mathcal{C}')\delta'/4 > \varepsilon(\mathcal{C})\delta'/8 = (m + 1)\delta'/8 \geq m + 1.$$

On the other hand,  $\bar{\varepsilon}(\mathcal{C}') \leq m$ . Thus,  $L(G)$  contains  $\mathcal{C}_m$ , a contradiction.

*Case 1.2:*  $\mathcal{C}$  has no chord. Since  $\delta' \geq 8$ ,  $\mathcal{C}$  cannot be a Hamilton cycle of  $G$ . Let  $u$  be a vertex in  $V(G) \setminus V(\mathcal{C})$ . If  $u$  is adjacent to at least three vertices of  $\mathcal{C}$ , then  $G$  contains a cycle  $\mathcal{C}'$  with  $\varepsilon(\mathcal{C})/2 < \varepsilon(\mathcal{C}') \leq m$  and we obtain a contradiction as in Case 1.1. Hence,  $u$  is adjacent to at most two vertices of  $\mathcal{C}$ . Defining  $p$  as the number of edges incident with exactly one vertex of  $\mathcal{C}$ , we thus have

$$p \leq 2|V(G) \setminus V(\mathcal{C})| = 2(n - m - 1).$$

On the other hand, since  $G$  has no chords,

$$p = \sum_{e \in E(\mathcal{C})} (d(e) - 4)/2 \geq (m + 1)(\delta' - 4)/2.$$

It follows that  $(m + 1)(\delta' - 4)/2 \leq 2(n - m - 1)$  or, equivalently,

$$m \leq \frac{4n - \delta'}{\delta'}.$$

*Case 2:*  $r \geq 2$ . Let  $H$  be the graph with  $V(H) = \{D_1, \dots, D_r\}$  and  $D_i D_j \in E(H)$  iff  $V(D_i) \cap V(D_j) \neq \emptyset$  ( $i \neq j$ ). Since  $H$  is connected, at least two vertices of  $H$  are not cut vertices of  $H$ . Equivalently, there are at least two values of  $j$  for which  $\bigcup_{1 \leq i \leq r, i \neq j} D_i$  is a connected subgraph of  $G$ , and hence a circuit of  $G$ .

Assume, without loss of generality, that  $\mathcal{C}' = \bigcup_{i=2}^r D_i$  and  $\mathcal{C}'' = D_1 \cup \bigcup_{i=3}^r D_i$  are circuits of  $G$ . We have

$$\begin{aligned} \bar{\varepsilon}(\mathcal{C}') &\geq |E(\mathcal{C}')| + |E(D_1) \cap \bar{E}(\mathcal{C}')| + |E(D_2 - V(\mathcal{C}''))|(\delta' - 4)/4 \\ &= |E(\mathcal{C})| - |E(D_1 - V(\mathcal{C}'))| + |E(D_2 - V(\mathcal{C}''))|(\delta' - 4)/4. \end{aligned}$$

On the other hand, since  $L(G)$  does not contain  $\mathcal{C}_m$ ,

$$\tilde{e}(\mathcal{C}') \leq m - 1 = e(\mathcal{C}) - 2.$$

It follows that  $|E(D_1 - V(\mathcal{C}'))| \geq |E(D_2 - V(\mathcal{C}''))|(\delta' - 4)/4 + 2$  and hence, since  $\delta' \geq 8$ ,

$$|E(D_1 - V(\mathcal{C}'))| > |E(D_2 - V(\mathcal{C}''))|.$$

But then by symmetry we also have

$$|E(D_2 - V(\mathcal{C}''))| > |E(D_1 - V(\mathcal{C}'))|.$$

This contradiction completes the proof.  $\square$

In a way similar to the proof of Lemma 5, one easily proves the following.

**Lemma 5'.** *Let  $G$  be a graph of order  $n$ , girth at least 5, and with minimum edge degree  $\delta' \geq 8$  such that  $L(G)$  contains  $\mathcal{C}_{m+1}$  but not  $\mathcal{C}_m$ . Then*

$$m \leq \frac{2n - \delta' + 2}{\delta' - 2}.$$

The proofs of our main results also need the following lemmas.

**Lemma 6.** *Let  $G$  be a graph of order  $n$  with  $\delta' > \sqrt{2n+1}$  ( $n \geq 72$ ). Then  $g(G) \leq \Delta(G) + 1$ .*

**Proof.** Since  $n \geq 72$ , we have  $\delta' \geq 13$ . Let  $G$  satisfy the hypothesis of the lemma and  $\mathcal{C}$  be a shortest cycle of  $G$ .

We distinguish the following two cases.

*Case 1:*  $\delta(G) = 1$ . Obviously,  $\Delta(G) \geq \delta' - 1$ . Set  $E_1(\mathcal{C}) = \{e = uv \in E(\mathcal{C}) \mid \text{either } u \text{ or } v \text{ is adjacent to an isolated vertex in } G - V(\mathcal{C})\}$ ,  $E_2(\mathcal{C}) = E(\mathcal{C}) \setminus E_1(\mathcal{C})$  and  $V_2(\mathcal{C}) = V(G[E_2(\mathcal{C})])$ .

We have the following claim.

$$d(e) \geq (\delta' - 1) + 2 = \delta' + 1 \quad \text{for any edge } e \in E_1(\mathcal{C}). \quad (1)$$

Assuming  $e(\mathcal{C}) \geq \Delta(\mathcal{C}) + 2$ , we thus have  $e(\mathcal{C}) \geq \delta' + 1 \geq 14$ , implying that

$$|\{uv \in E(G) : v \in V(\mathcal{C})\}| \leq 1 \quad \text{for any } u \in V(G) \setminus V(\mathcal{C}). \quad (2)$$

By the choice of  $\mathcal{C}$ , we obtain that for any vertex  $x \in V_2(\mathcal{C})$ , there exists a set  $\{p_i\}$  of  $d(x) - 2$  paths of length 2 such that  $V(p_i) \cap V(\mathcal{C}) = \{x\}$  and  $V(p_i) \cap V(p_j) = \{x\}$  ( $i \neq j$ ).

Defining  $P_1$  as the set of such paths, we thus obtain that any pair of paths in  $P_1$  have at most one common vertex in  $G$  and

$$|P_1| \geq \sum_{e \in E_2(\mathcal{C})} (d(e) - 4)/2 \geq |E_2(\mathcal{C})|/2. \quad (3)$$

Using (1)–(3) and  $\delta' > \sqrt{2n+1}$ , we obtain

$$\begin{aligned} n &\geq \left( \sum_{e \in E(\mathcal{C})} (d(e) - 4) \right) / 2 + \varepsilon(\mathcal{C}) + |P_1| \\ &\geq \left( \sum_{e \in E_1(\mathcal{C})} d(e) + \sum_{e \in E_2(\mathcal{C})} d(e) \right) / 2 + |P_1| - \varepsilon(\mathcal{C}) \\ &\geq ((\delta' + 1)|E_1(\mathcal{C})| + \delta'|E_2(\mathcal{C})|) / 2 + |E_2(\mathcal{C})| / 2 - \varepsilon(\mathcal{C}) \\ &= (\delta' + 1)(|E_1(\mathcal{C})| + |E_2(\mathcal{C})|) / 2 - \varepsilon(\mathcal{C}) \\ &= \varepsilon(\mathcal{C})(\delta' - 1) / 2 \geq (\delta' + 1)(\delta' - 1) / 2 > n, \end{aligned}$$

a contradiction.

*Case 2:*  $\delta(G) > 1$ . Assuming  $\varepsilon(\mathcal{C}) \geq \Delta(G) + 2$ , we thus have  $\varepsilon(\mathcal{C}) \geq (\delta' + 4) / 2 \geq \frac{17}{2}$ , implying that  $\varepsilon(\mathcal{C}) \geq 9$ . By the choice of  $\mathcal{C}$  and  $\delta(G) > 1$ , we have that for any vertex  $x \in V(\mathcal{C})$ , there exists a set  $\{p_i\}$  of  $d(x) - 2$  paths of length 2 such that  $V(p_i) \cap V(\mathcal{C}) = \{x\}$  and  $V(p_i) \cap V(p_j) = \{x\}$  ( $i \neq j$ ).

Defining  $P_2$  as the set of such paths, we thus obtain that any pair of paths in  $P_2$  have at most one common vertex in  $G$  and

$$\begin{aligned} |P_2| &\geq \sum_{u \in V(\mathcal{C})} (d(u) - 2), \text{ implying that} \\ n &\geq 2 \sum_{u \in V(\mathcal{C})} (d(u) - 2) + \varepsilon(\mathcal{C}) = 2 \left( \sum_{e \in E_1(\mathcal{C})} (d(e) - 4) \right) / 2 + \varepsilon(\mathcal{C}) \\ &\geq ((\delta' - 4) + 1)\varepsilon(\mathcal{C}) \geq (\delta' - 3)(\delta' + 4) / 2 > n, \end{aligned}$$

a contradiction. This completes the proof.  $\square$

**Lemma 7.** Let  $G$  be a graph of order  $n$ , girth at least 5, and with  $g(G)(\delta' - 2)^2 + 2(\delta' - 2) \geq 4n$ . Then

$$\lambda(L(G)) = [g(G), \text{cr}(L(G))].$$

**Proof.** Let  $G$  satisfy the hypothesis of the lemma and  $\mathcal{C}$  be a shortest cycle of  $G$ , i.e.,  $\varepsilon(\mathcal{C}) = g(G)$ . We have

$$\begin{aligned} \bar{\varepsilon}(\mathcal{C}) &= \left( \sum_{e \in E(\mathcal{C})} (d(e) - 4) \right) / 2 + \varepsilon(\mathcal{C}) \\ &= \left( \sum_{e \in E(\mathcal{C})} d(e) \right) / 2 - \varepsilon(\mathcal{C}) \\ &\geq (\delta' - 2)\varepsilon(\mathcal{C}) / 2 \\ &= g(G)(\delta' - 2) / 2. \end{aligned} \tag{4}$$

By  $g(G)(\delta' - 2)^2 + 2(\delta' - 2) \geq 4n$ ,

$$g(G)(\delta' - 2)/2 \geq 2n/(\delta' - 2) - 1. \quad (5)$$

Using (4) and (5), Lemma 5' and Theorem 4, we obtain  $\lambda(L(G)) = [g(G), \text{cr}(L(G))]$ , which completes the proof of Lemma 7.  $\square$

**Proof of Theorem 1.** Let  $G$  satisfy the hypothesis of Theorem 1 and  $\mathcal{C}$  be a shortest cycle of  $G$ , i.e.,  $\varepsilon(\mathcal{C}) = g(G)$ . From Lemmas 6, 7 and Theorem 4, we have the  $L(G)$  is subpancyclic.

Next, we construct a family of graphs of order  $n$  with  $\delta' = \sqrt{2n+1}$  such that their line graphs are Hamiltonian but not pancyclic.

For any integer  $d = 2k+1 \geq 3$ , define the graph  $G_d$  as follows. Let  $\mathcal{C} = u_1 u_2 \dots u_{d+1} u_1$  be a cycle of length  $d+1$  and let  $H_1, H_2, \dots, H_{k+1}$  be  $k+1$  copies of the empty graph of order  $d-3$  such that  $\mathcal{C}, H_1, H_2, \dots, H_{k+1}$  are pairwise disjoint.

Now,  $G_d$  is obtained from  $\mathcal{C} \cup \bigcup_{i=1}^{k+1} H_i$  by joining each vertex of  $H_i$  to  $u_{2i}$ , for  $i = 1, 2, \dots, k+1$ . We have

$$\delta'(G_d) = d \quad (6)$$

and

$$|V(G_d)| = (d+1) + (k+1)(d-3) = (d^2 - 1)/2. \quad (7)$$

Clearly,  $\mathcal{C}$  is a cycle of  $G$  with  $\bar{\varepsilon}(\mathcal{C}) = |E(G)|$ , hence

$$L(G_d) \text{ is Hamiltonian.} \quad (8)$$

Obviously,  $L(G_d)$  does not contain  $C_d$  and hence

$$L(G_d) \text{ is not pancyclic.} \quad (9)$$

Using (6)–(9), we obtain that Theorem 1 is best possible in the sense that the condition  $\delta' > \sqrt{2n+1}$  cannot be relaxed, even under the condition that  $L(G)$  is Hamiltonian. This completes the proof of Theorem 1.  $\square$

**Proof of Theorem 2.** Let  $G$  satisfy the hypothesis of the theorem. Assume  $L(G)$  is not subpancyclic. Set  $m = \max\{i \leq \text{cr}(L(G)) \mid L(G) \text{ does not contain } \mathcal{C}_i\}$ . Then  $m \leq \text{cr}(L(G)) - 1$  and  $L(G)$  contains  $C_{m+1}$ . By Theorem 4,  $G$  contains a circuit  $\mathcal{C}$  with  $\varepsilon(\mathcal{C}) \leq m+1 < \bar{\varepsilon}(\mathcal{C})$ . In fact,  $\varepsilon(\mathcal{C}) = m+1$ , otherwise  $L(G)$  contains  $\mathcal{C}_m$ . Using arguments similar to those in the proof of Lemma 5, we obtain

**Claim 1.**  $\mathcal{C}$  is a cycle without any chord and  $\mathcal{C}$  cannot be a Hamiltonian cycle of  $G$ .

Let  $w$  be a vertex in  $V(G) \setminus V(\mathcal{C})$ . It is easy to see that

**Claim 2.**  $w$  is adjacent to at most two vertices of  $\mathcal{C}$  and the distance (in  $\mathcal{C}$ ) between any pair of  $N(w) \cap V(\mathcal{C})$  is 2.

Otherwise  $G$  contains a cycle  $\mathcal{C}'$  with  $\varepsilon(\mathcal{C})/2 < \varepsilon(\mathcal{C}') \leq m$ , and we obtain a contradiction as in Case 1.1 in Lemma 5.

By Theorem 1, we only need to consider the case that  $g(G) = 4$ . Let  $\mathcal{C}_4$  be 4-cycle of  $G$ . By the arguments similar to the proof of Lemma 7, we have  $\bar{\varepsilon}(\mathcal{C}_4) \geq 2(\delta' - 2)$ . By Theorem 4,  $\lambda(L(G)) \geq [3, 2\delta' - 4]$ . Clearly,

$$\varepsilon(\mathcal{C}) \geq (2\delta' - 4) + 2 = 2\delta' - 2. \quad (10)$$

Using (10), Claims 1, 2 and  $\delta'^2 - \delta' > 2n$ , we obtain

$$\begin{aligned} n &\geq \left( \sum_{e \in E(\mathcal{C})} (d(e) - 4) \right) / 4 + \varepsilon(\mathcal{C}) \\ &= \left( \sum_{e \in E(\mathcal{C})} d(e) \right) / 4 \geq \varepsilon(\mathcal{C})\delta' / 4 \\ &\geq 2(\delta' - 1)\delta' / 4 > n, \end{aligned}$$

a contradiction, which implies that  $L(G)$  is subpancyclic.

Next, we construct a family of graphs of order  $n$  with  $\delta'^2(G) - \delta'(G) = 2n$  such that their line graphs are Hamiltonian but not pancyclic.

For any integer  $d = 2k + 1 \geq 3$ , define the graph  $G_d$  as follows.

Let  $\mathcal{C} = u_1, u_2, \dots, u_{4k}, u_1$  be a cycle of length  $4k$  and let  $H_1, H_2, \dots, H_k$  be  $k$  copies of  $2k - 3$  isolated vertices such that  $\mathcal{C}, H_1, H_2, \dots, H_k$  are pairwise disjoint. Now  $G_d$  is obtained from  $\mathcal{C} \cup \bigcup_{i=1}^k H_i$  by joining each vertex of  $H_i$  to  $u_{4i-3}$  and  $u_{4i-1}$ , for  $i = 1, 2, \dots, k$ . We have

$$\delta'(G) = d \quad (11)$$

and

$$|V(G_d)| = 4k + k(d - 4) = (d^2 - d)/2. \quad (12)$$

Clearly,  $\mathcal{C}$  is a cycle of  $G$  with  $\bar{\varepsilon}(\mathcal{C}) = |E(G)|$ , hence

$$L(G_d) \text{ is Hamiltonian.} \quad (13)$$

Obviously,  $L(G_d)$  does not contain  $\mathcal{C}_{4k-1} = \mathcal{C}_{2d-3}$  and hence

$$L(G_d) \text{ is not pancyclic.} \quad (14)$$

Using (11)–(14), we obtain that Theorem 2 is best possible in the sense that the condition  $\delta'^2 - \delta' > 2n$  cannot be relaxed, even under the condition that  $L(G)$  is Hamiltonian. This completes the proof of Theorem 2.  $\square$

From the proofs of Theorems 1 and 2, two important corollaries can easily be obtained.

**Corollary 8.** *Let  $G$  be a graph of order  $n \geq 72$ , girth at least 5, and with  $\delta' > \sqrt{2n+1}$ . Then the line graph  $L(G)$  of  $G$  is pancyclic whenever  $L(G)$  is Hamiltonian, and the result is best possible.*

**Corollary 9.** *Let  $G$  be a graph of order  $n \geq 72$ , girth at least 4, and with  $\delta'^2 - \delta' > 2n$ . Then the line graph  $L(G)$  of  $G$  is pancyclic whenever  $L(G)$  is Hamiltonian, and the result is best possible.*

In fact, Corollary 9 verifies the following conjecture in [5] for those graphs of order  $n \geq 72$  with girth at least 4.

**Conjecture 10** (Liming Xiong [5]). *Let  $G$  be a graph with  $\delta'^2 - \delta' > 2n$ . Then the line graph  $L(G)$  of  $G$  is pancyclic whenever  $L(G)$  is Hamiltonian, unless  $G$  is isomorphic to  $C_4$ ,  $C_5$ , or the Petersen graph.*

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